

Section 7.5

Equivalence Relations

Now we group properties of relations together to define new types of important relations.

Definition: A relation R on a set A is an *equivalence relation* iff R is

- reflexive
- symmetric

and

- transitive
-

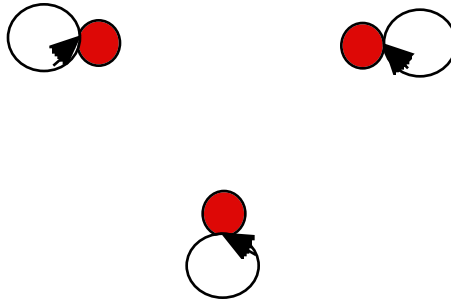
It is easy to recognize equivalence relations using digraphs.

- The subset of all elements related to a particular element forms a universal relation (contains all possible arcs) on that subset. The (sub)digraph representing the subset is called a *complete* (sub)digraph. All arcs are present.

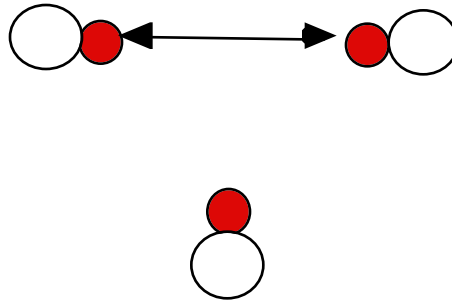
- The number of such subsets is called the *rank* of the equivalence relation

Examples:

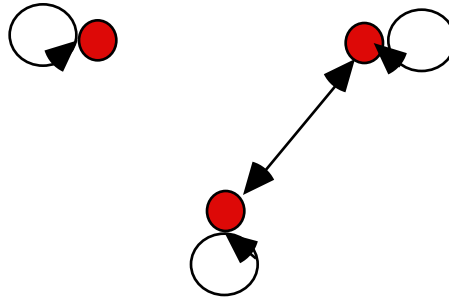
A has 3 elements:



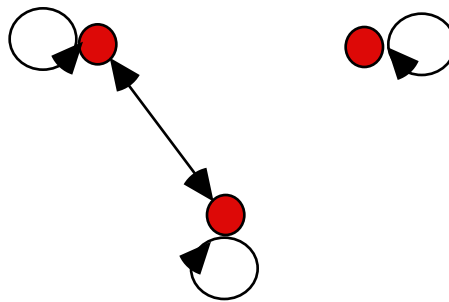
rank = 3



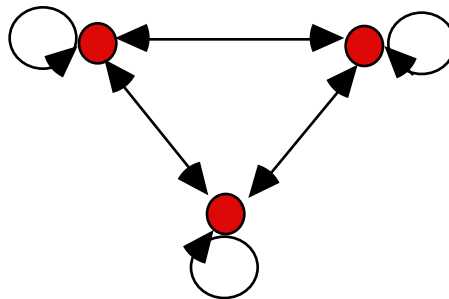
rank = 2



rank = 2



rank = 2



rank = 1

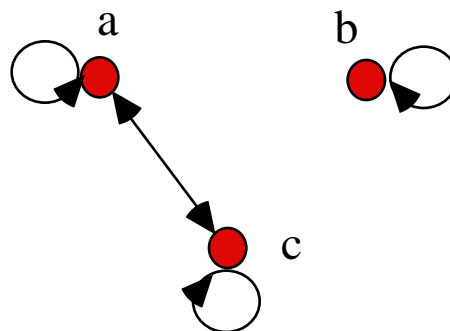
-
- Each of the subsets is called an *equivalence class*.

- A bracket around an element means the equivalence class in which the element lies.

$$[x] = \{y \mid \langle x, y \rangle \text{ is in } R\}$$

- The element in the bracket is called a *representative* of the equivalence class. We could have chosen any one.
-

Examples:



$$[a] = \{a, c\}, [c] = \{a, c\}, [b] = \{b\}.$$

$$\text{rank} = 2$$

An interesting counting problem:

Count the number of equivalence relations on a set A with n elements. Can you find a recurrence relation?

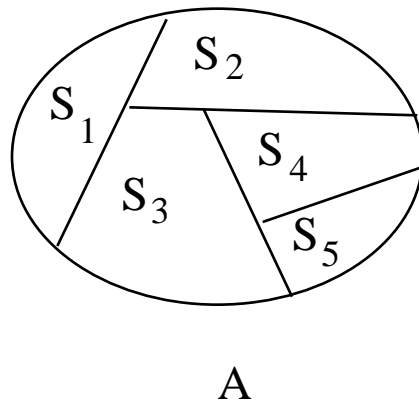
The answers are

- 1 for $n = 1$
- 3 for $n = 2$
- 5 for $n = 3$

How many for $n = 4$?

Definition: Let S_1, S_2, \dots, S_n be a collection of subsets of A . Then the collection forms a *partition* of A if the subsets are nonempty, disjoint and *exhaust* A :

- S_i
- $S_i \cap S_j = \emptyset$ if $i \neq j$
- $\bigcup S_i = A$



Theorem: The equivalence classes of an equivalence relation R *partition* the set A into disjoint nonempty subsets whose union is the entire set.

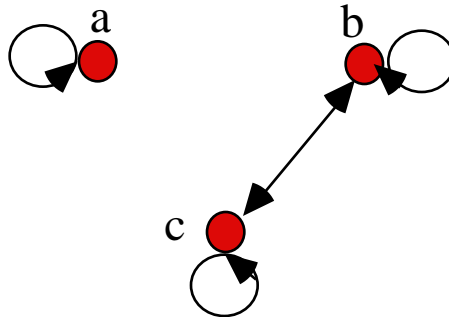
This partition is denoted A/R and called

- the *quotient set*, or
- *the partition of A induced by R* , or,
- *A modulo R .*

no. of partitions: $B_{n+1} = \sum_{k=0}^n B_n * C(n,k)$

Examples:

- *Ex. 1, p. 508*
- *Ex. 4, p. 509*
- *Ex. 9, p. 512, problem 18, p. 514*



$$A = [a] \quad [b] = [a] \quad [c] = \{a\} \quad \{b, c\}$$

$$\text{rank} = 2$$

Theorem: Let R be an equivalence relation on A . Then either

$$[a] = [b]$$

or

$$[a] \cap [b] = \emptyset$$

Theorem: If R_1 and R_2 are equivalence relations on A then $R_1 \cap R_2$ is an equivalence relation on A .

Proof: It suffices to show that the intersection of

- reflexive relations is reflexive,

- symmetric relations is symmetric,

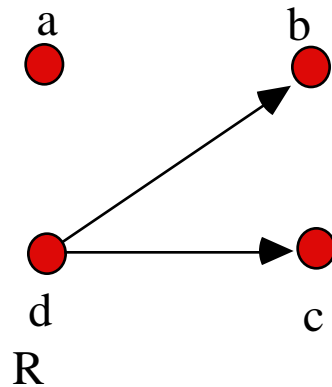
and

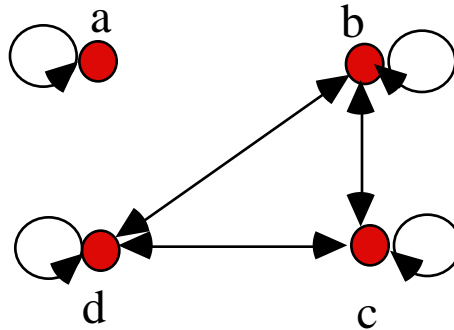
- transitive relations is transitive.

You provide the details.

Definition: Let R be a relation on A . Then the reflexive, symmetric, transitive closure of R , $\text{tsr}(R)$, is an equivalence relation on A , called the *equivalence relation induced by R* .

Example:





$$\begin{aligned} & \text{tsr}(R) \\ & \text{rank} = 2 \\ A = [a] \quad [b] = \{a\} \quad \{b, c, d\} \\ A/R = \{\{a\}, \{b, c, d\}\} \end{aligned}$$

Theorem: $\text{tsr}(R)$ is an equivalence relation

Proof:

We have to be careful and show that $\text{tsr}(R)$ is still symmetric and reflexive.

- Since we only add arcs vs. deleting arcs when computing closures it must be that $\text{tsr}(R)$ is reflexive since all loops $\langle x, x \rangle$ on the diagraph must be present when constructing $r(R)$.

- If there is an arc $\langle x, y \rangle$ then the symmetric closure of $r(R)$ ensures there is an arc $\langle y, x \rangle$.

• Now argue that if we construct the transitive closure of $sr(R)$ and we add an edge $\langle x, z \rangle$ because there is a path from x to z , then there must also exist a path from z to x (why?) and hence we also must add an edge $\langle z, x \rangle$. Hence the transitive closure of $sr(R)$ is symmetric.

Q. E. D.
