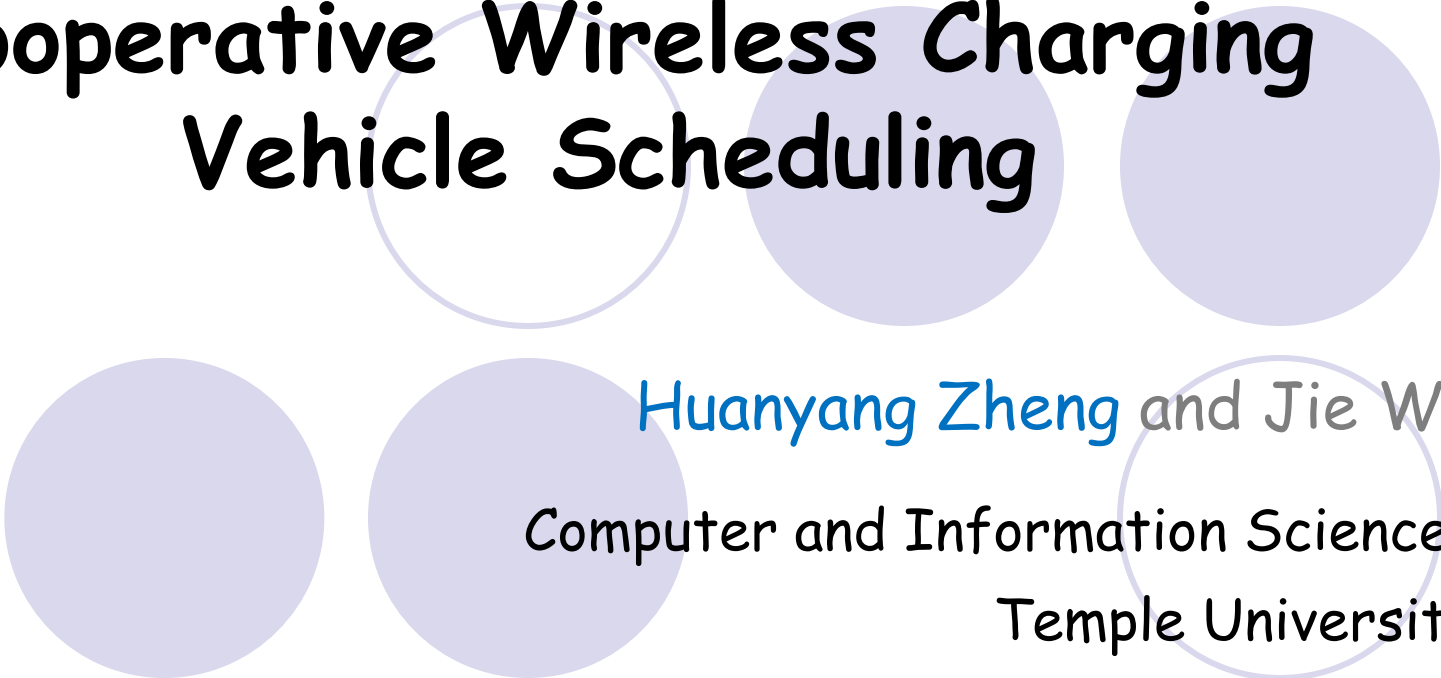


# Cooperative Wireless Charging Vehicle Scheduling



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# 1. Introduction



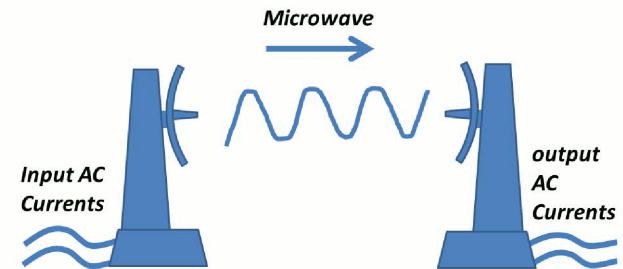
- Limited lifetime of battery-powered WSNs
- Possible solutions
  - Energy conservation
    - Cannot compensate for energy depletion
  - Energy harvesting (or scavenging)
    - Unstable, unpredictable, uncontrollable...
  - Sensor reclamation
    - Costly, impractical (deep ocean, bridge surface ...)

(WSNs: Wireless Sensor Networks)

## 2. State of the Art

- The enabling technology

- Wireless energy transfer (Kurs '07)
- Wireless Power Consortium



- Wireless charging vehicles (WCV)

- A WCV moves from one location to another for wireless charging
- Extended from **mobile sink** in WSNs and **ferry** in DTNs
- Energy consumption

- The movement of WCV
- The energy charging process

(DTNs: Delay Tolerant Networks)

# 3. Collaborative Coverage & Charging

- Most existing methods

- A WCV is fast enough to charge all sensors in a cycle
- A WCV has sufficient energy to replenish an entire WSN (and return to the BS)

- Collaborative approach using multiple WCVs

**Problem:** Determine the minimum number of WCVs (with an unrestricted capacity but limitations on speed) to cover sensors with uniform/non-uniform recharge frequencies

# Problem Description

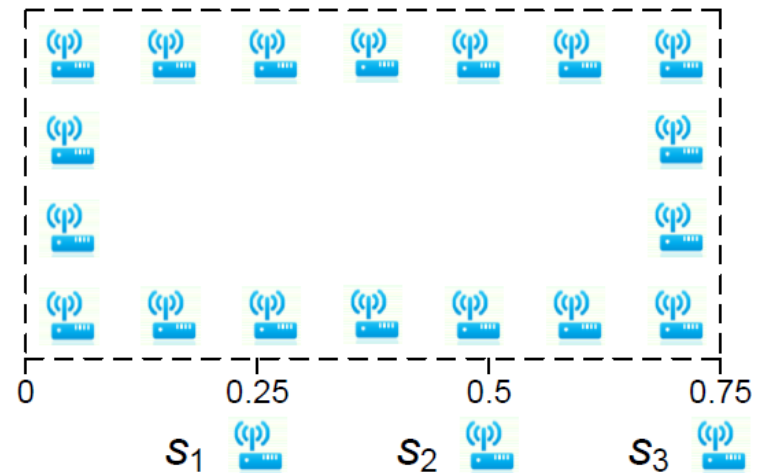
- A toy example

- A rectangle track with a circumference of 3.75 is densely covered with sensors with a recharge frequency of  $f=1$  (WCV's max speed is 1)

- Sensors with  $f=2$  at 0.25 and 0.75

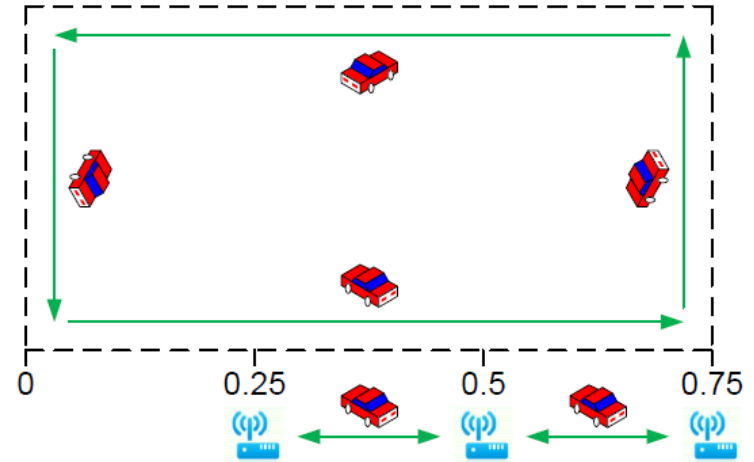
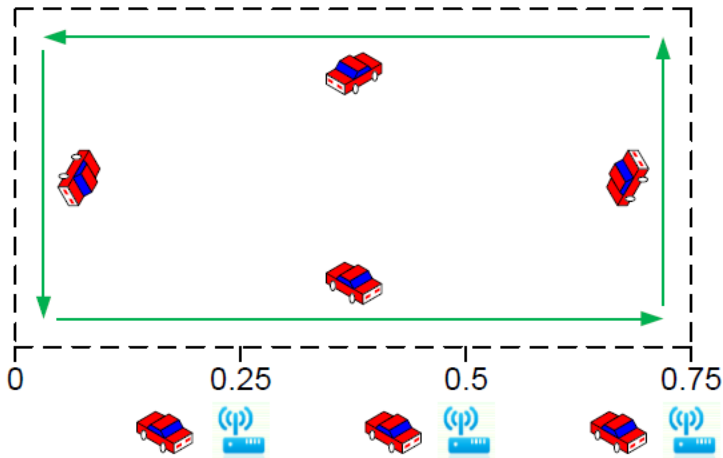
- A sensor with  $f=4$  at 0.5

- What is the minimum number of WCVs and the optimal trajectory of these MCs?

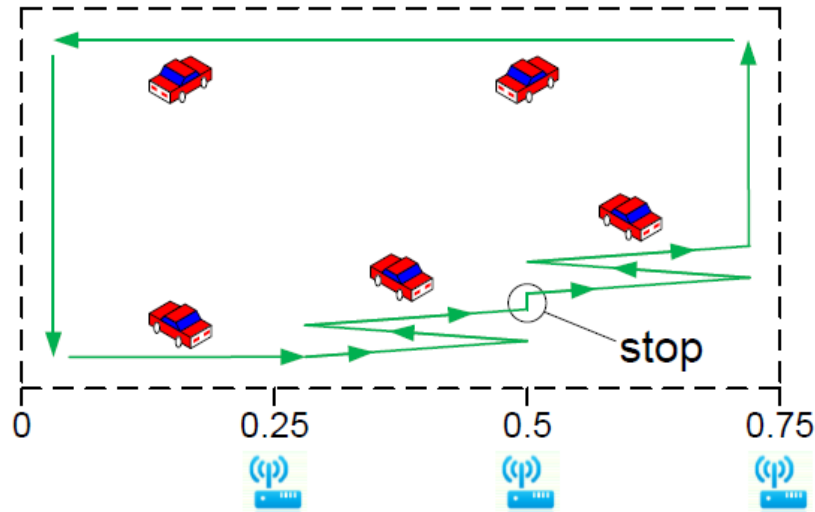


# Possible Solutions

- Assigning cars for sensors with  $f > 1$



- Optimal solution



# 4. Algorithm Design



- Line space

- Uniform frequency: optimal

- Non-uniform frequency: bound of 2

- Circle space

- Uniform frequency: optimal

- Non-uniform frequency: bound of 4

- Metric space

- Uniform frequency: bound of 2.5

- Non-uniform frequency: bound of  $5\log_2 f_{\max}/f_{\min}$

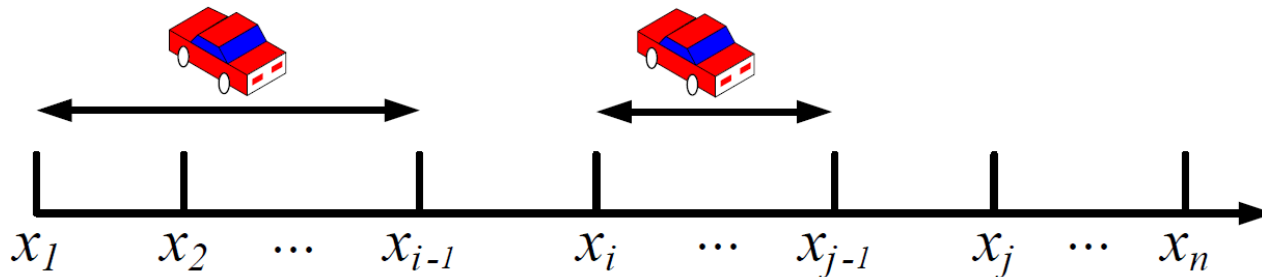
# Line space (non-uniform frequency)

- Back and Forth (BF)

$\text{BF}(x_1, \dots, x_n; f_1, \dots, f_n)$ :

When  $n \neq 0$ , generate a WCV that goes back and forth as far as possible at its full speed (covering  $x_1, \dots, x_{i-1}$ );

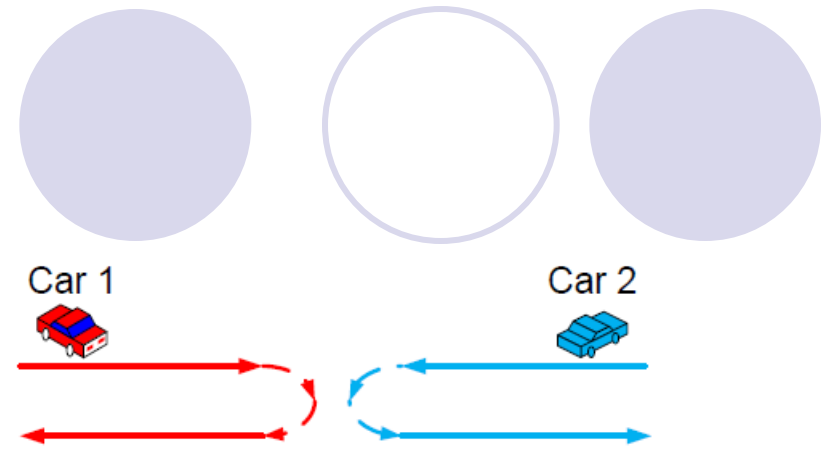
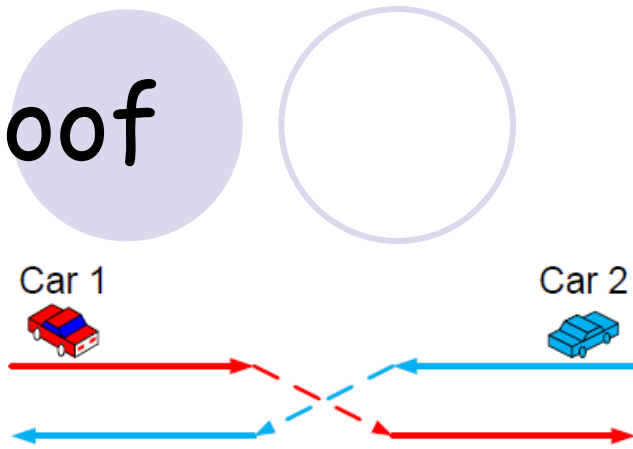
$\text{BF}(x_i, \dots, x_n; f_i, \dots, f_n)$



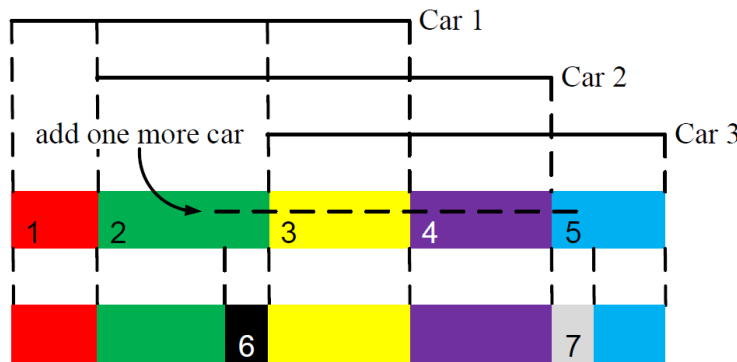
**Theorem 1:** BF guarantees an approximation ratio of 2 for its optimal solution

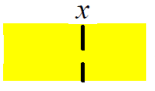


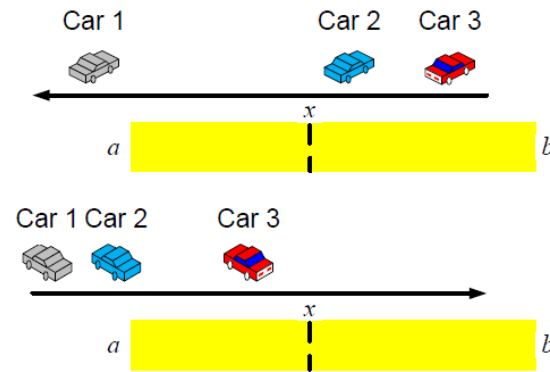
# Proof



- Two cars never meet or pass each other
- Partition the line into  $2k-1$  sub-regions based on the different car coverage ( $k$  is the optimal number of cars)
- Each sub-region can be served by one car going full speed
- One extra car is used when a circle is broken into a line



- 1-red  
2-green  
3-yellow
- $a$    $b$
- 4-purple 5-blue  
6-black 7-gray



$$2(x-a) \leq f_x \text{ and } 2(b-x) \leq f_x$$



# Line space (uniform frequency)

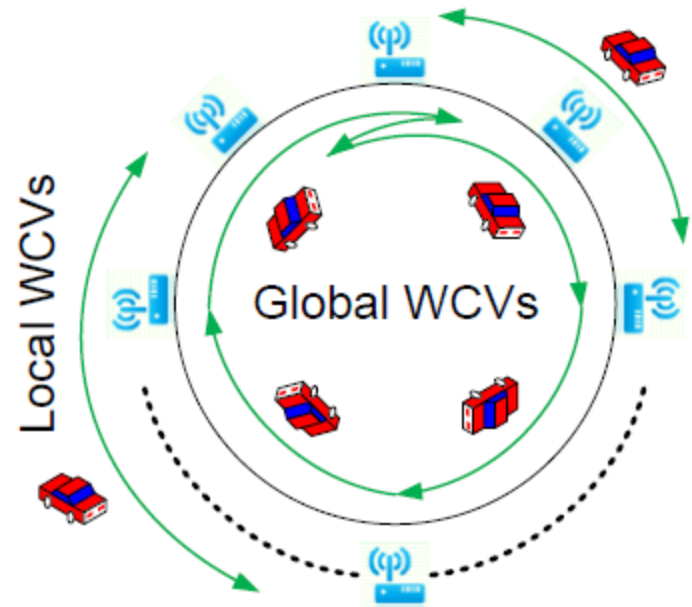
- The greedy solution becomes optimal
  - Proof by contradiction

# Cycle space (non-uniform frequency)

- Break the cycle space into line space
  - Use the greedy solution on the line space
  - An approximation ratio of 4

# Cycle space (uniform frequency)

- $M_1$ : There are  $C_1$  WCVs moving continuously around the circle
- $M_2$ : There are  $C_2$  WCVs moving inside the fixed interval of length  $\frac{1}{2}$  so that all sensors are covered
- **Combined method**: it is either  $M_1$  or  $M_2$ ,  $C = \min \{C_1, C_2\}$ 
  - Optimal by contradiction
  - Otherwise, the WCVs meet



# Metric space (uniform frequency)

- Cycles in Minimum Forest (CMF)

$CMF(x_1, \dots, x_n; f_1, \dots, f_n)$ :

for a minimum spanning forest with  $k=1, 2, \dots$  and  $n-1$  edges

construct a TSP for each connected component of the minimum spanning forest and regard each TSP as a cycle space for scheduling WCVs

return the best one among all minimum spanning forests

**Theorem 2:** CMF guarantees an approximation ratio of 2.5 for the optimal solution

Proof idea: the total distances of WCVs in the optimal solution are larger than the minimum spanning forest

# Metric space (non-uniform frequency)

- Cover Sensors by Lifetimes (CSL)

$CSL(x_1, \dots, x_n; f_1, \dots, f_n)$ :

for  $i = 1$  to  $\log_2 f_{\max}/f_{\min}$

Sensors with  $2^{i-1} * f_{\min} < f < 2^i * f_{\min}$  are grouped

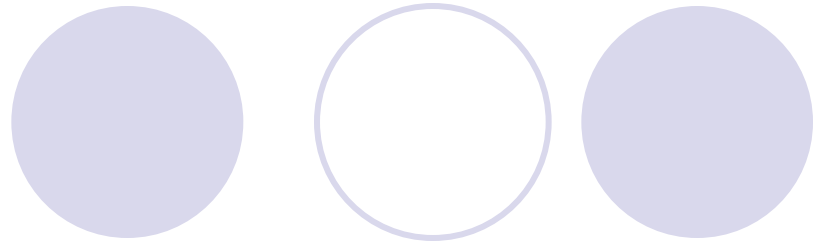
Schedule each sensor group with CMF

Return the best one among all minimum spanning forests

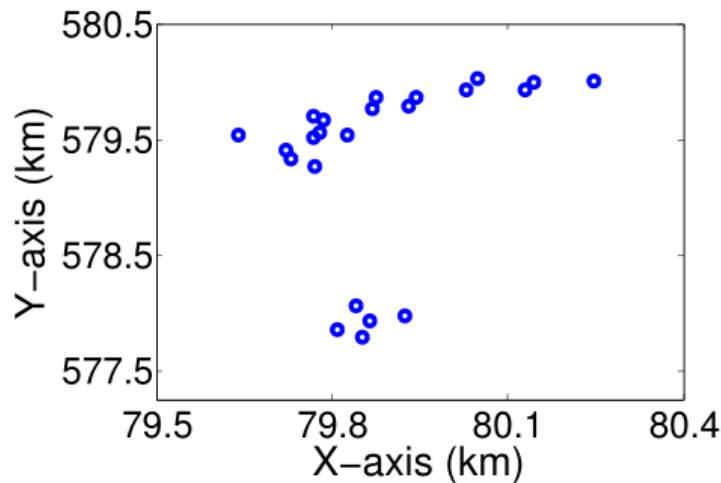
**Theorem 2:** CSL guarantees an approximation ratio of  $5 \log_2 f_{\max}/f_{\min}$  for the optimal solution

Proof idea: CMF has a ratio of 2.5; the number of sensor groups is  $\log_2 f_{\max}/f_{\min}$ ; each group has a ratio of 2

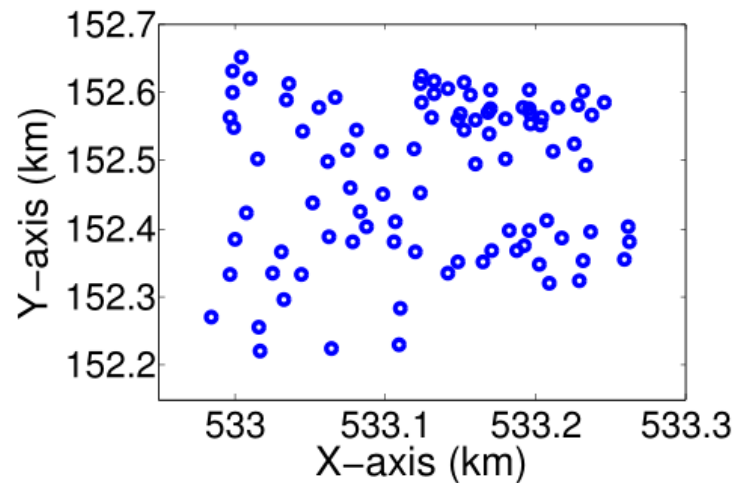
# 5. Experiments



- Two datasets (sensor locations)



(a) GSBD dataset.



(b) LUCE dataset.

- WCV speed is tuned from 10 km/h to 100 km/h
- Three sensor frequency (i.e., lifetime) distributions:
  - Uniform, normal, and exponential

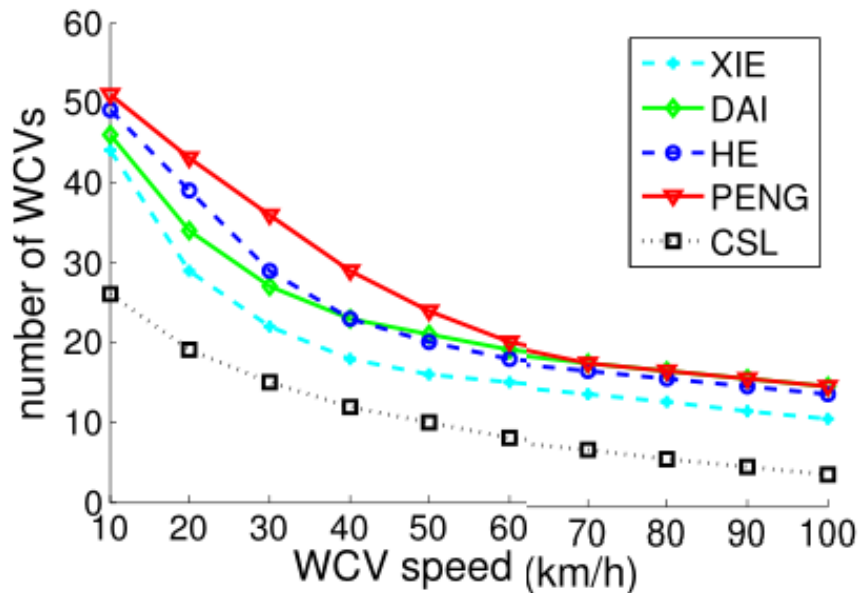
# Comparison Algorithms



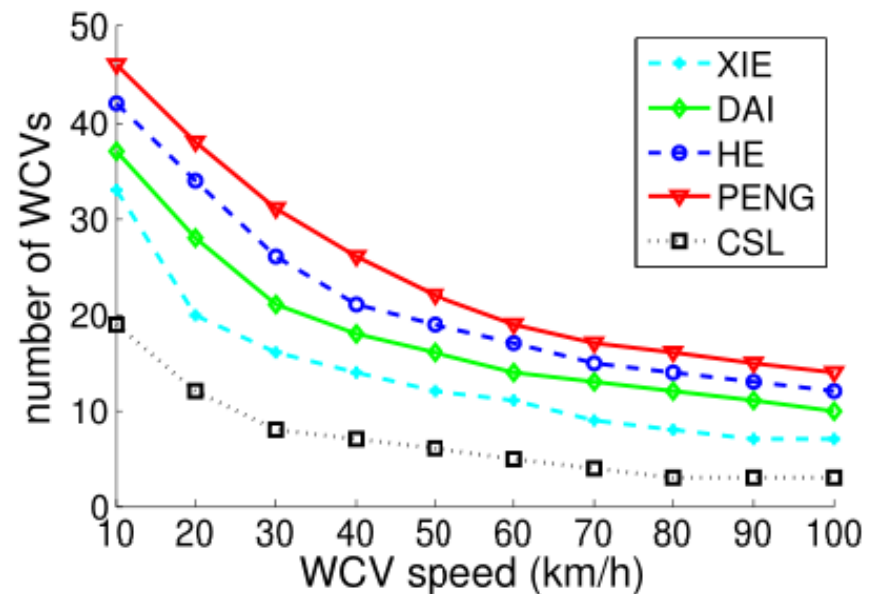
- **XIE** uses linear programming to near-optimally schedule only one WCV to recharge the sensors. It is iteratively applied until all sensors can run forever.
- **DAI** schedules WCVs according to vehicle routing problems. It can construct a forest among sensors, and assigns WCVs to each tree in the forest.
- **HE** is a on-demand approach in which each WCV prioritizes the nearest sensor that is close to its lifetime. It is also iteratively applied.
- **PENG** is a greedy approach that schedules WCVs according to sensor lifetimes. Sensors with shorter lifetimes are prioritized by WCVs.

# Experiment Results

- Impact of average sensor frequencies (GBSD)



(a) Uniform distribution,  $t \sim U(15, 25)$ .



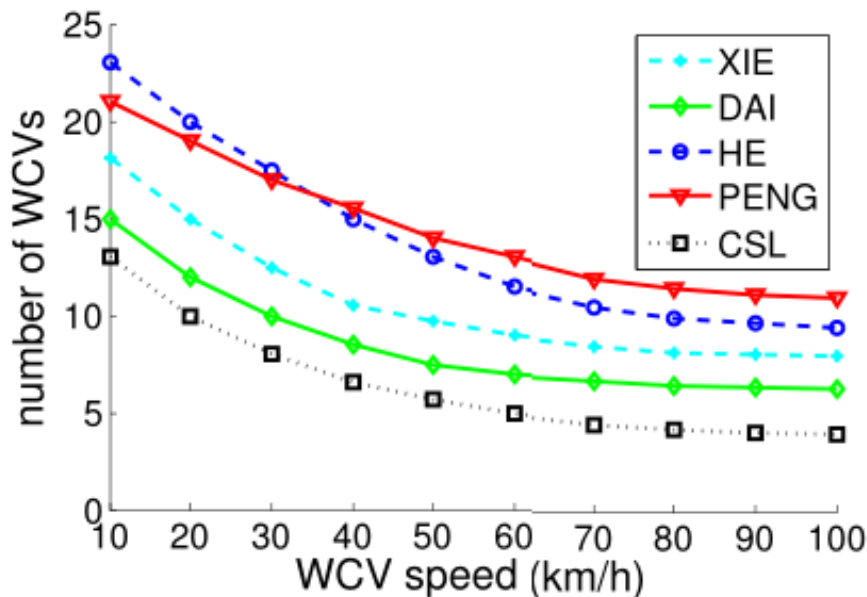
(b) Uniform distribution,  $t \sim U(25, 35)$ .

- More WCVs are needed for higher frequencies (smaller sensor lifetimes) and smaller WCV speeds

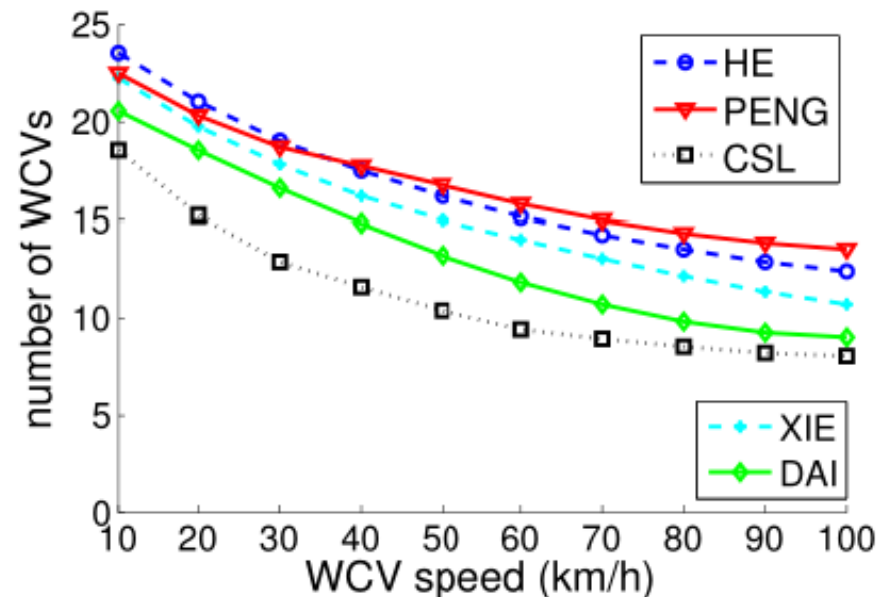


# Experiment Results

- Impact of sensor frequencies variance (GBSD)



(a) Uniform distribution,  $t \sim U(15, 25)$ .

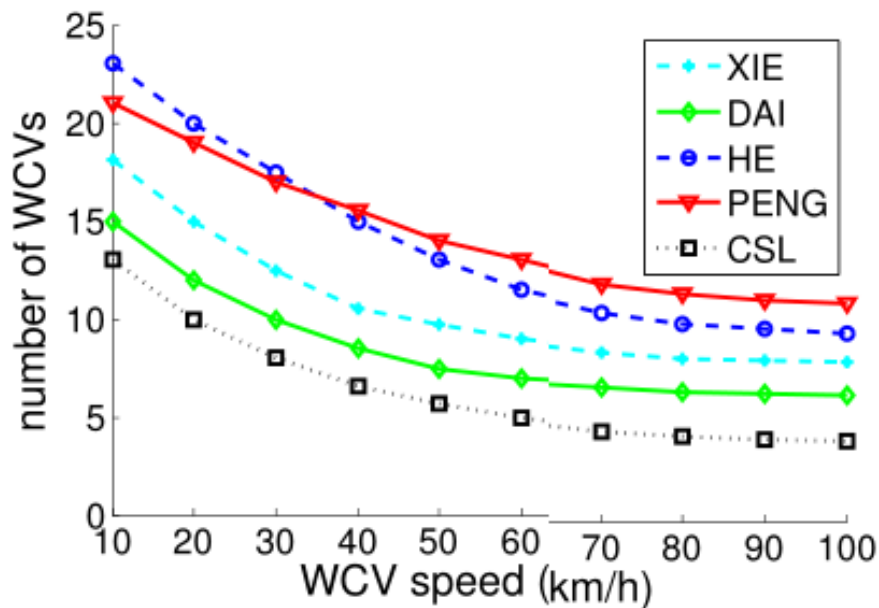


(b) Uniform distribution,  $t \sim U(5, 35)$ .

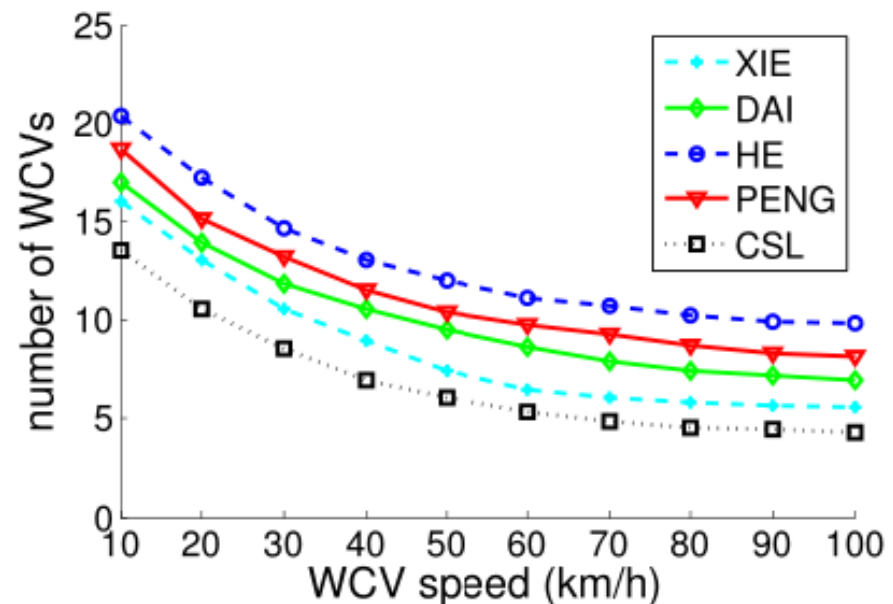
- More WCVs are needed for a higher frequency variance, especially when WCVs have high speeds

# Experiment Results

- Impact of frequency distribution difference (GBSD)



(a) Uniform distribution,  $t \sim U(15, 25)$ .



(b) Normal distribution,  $t \sim N(20, 5)$ .

- Not sensitive to distribution difference

# Experiment Summary



- Larger frequencies bring larger demands on WCVs
- Larger fluctuations of frequencies also bring larger demands on WCVs
- Not sensitive to frequency distribution differences

# 5. Conclusions



- Wireless energy transfer
- Collaborative mobile charging & coverage
  - Unlimited capacity, but limitations on speed
- Other extensions
  - Charging efficiency
  - WCVs as mobile sinks
  - ...